

A decorative vertical bar on the left side of the slide. It consists of a dark teal background with a white dotted vertical line running through its center. To the right of this bar, there are several orange circles of varying sizes, arranged in a cluster. The largest circle is at the top, with several smaller ones below and to its right. The entire slide is framed by thin orange vertical lines on the far left and far right.

PRINCIPLES OF OPERATING SYSTEMS

LECTURE 15 : SHARED PAGES AND SEGMENTATION

Shared Pages

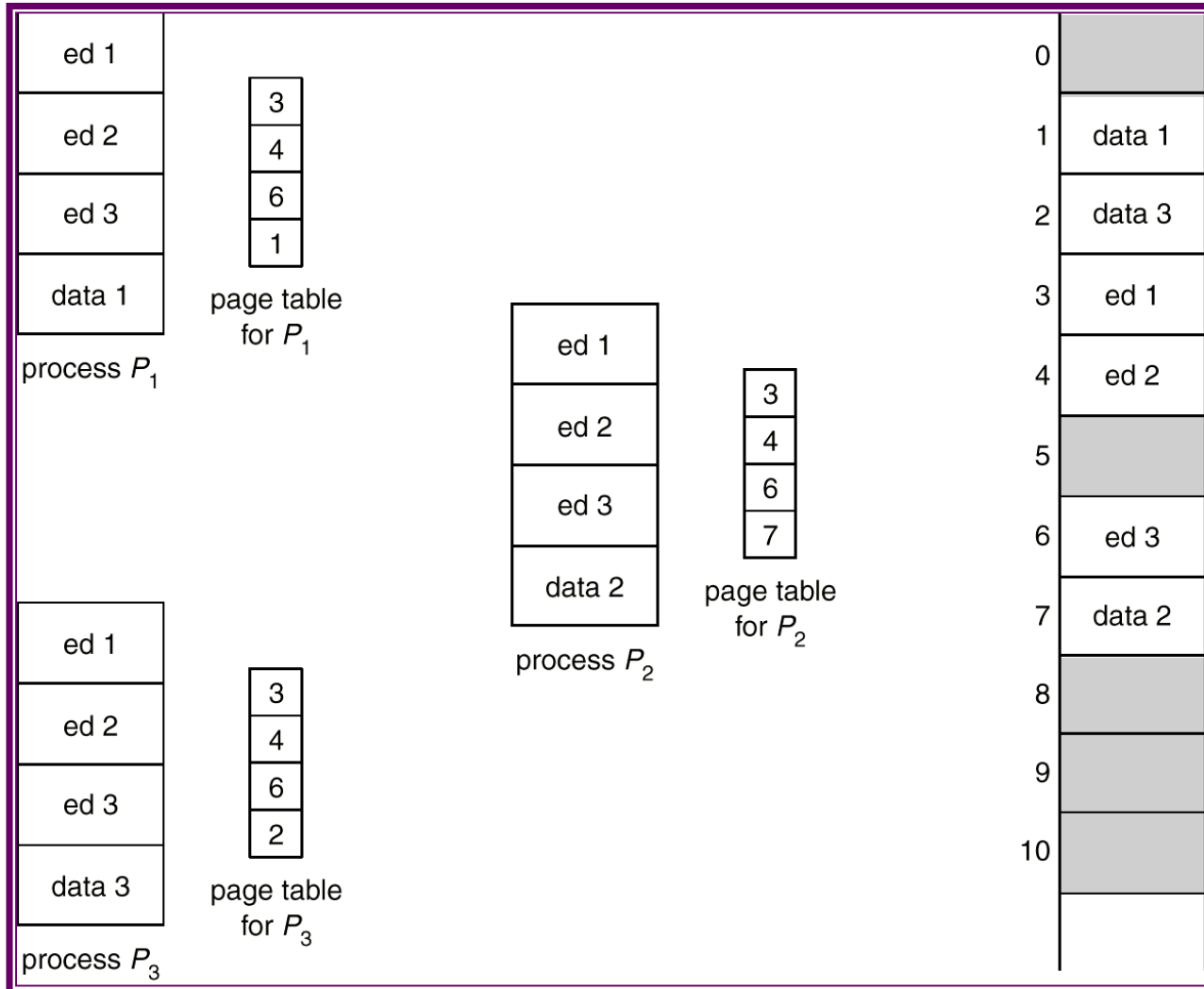
■ Shared code

- ☞ One copy of read-only (reentrant) code shared among processes (i.e., text editors, compilers, window systems).
- ☞ Shared code must appear in same location in the logical address space of all processes.

■ Private code and data

- ☞ Each process keeps a separate copy of the code and data.
- ☞ The pages for the private code and data can appear anywhere in the logical address space.

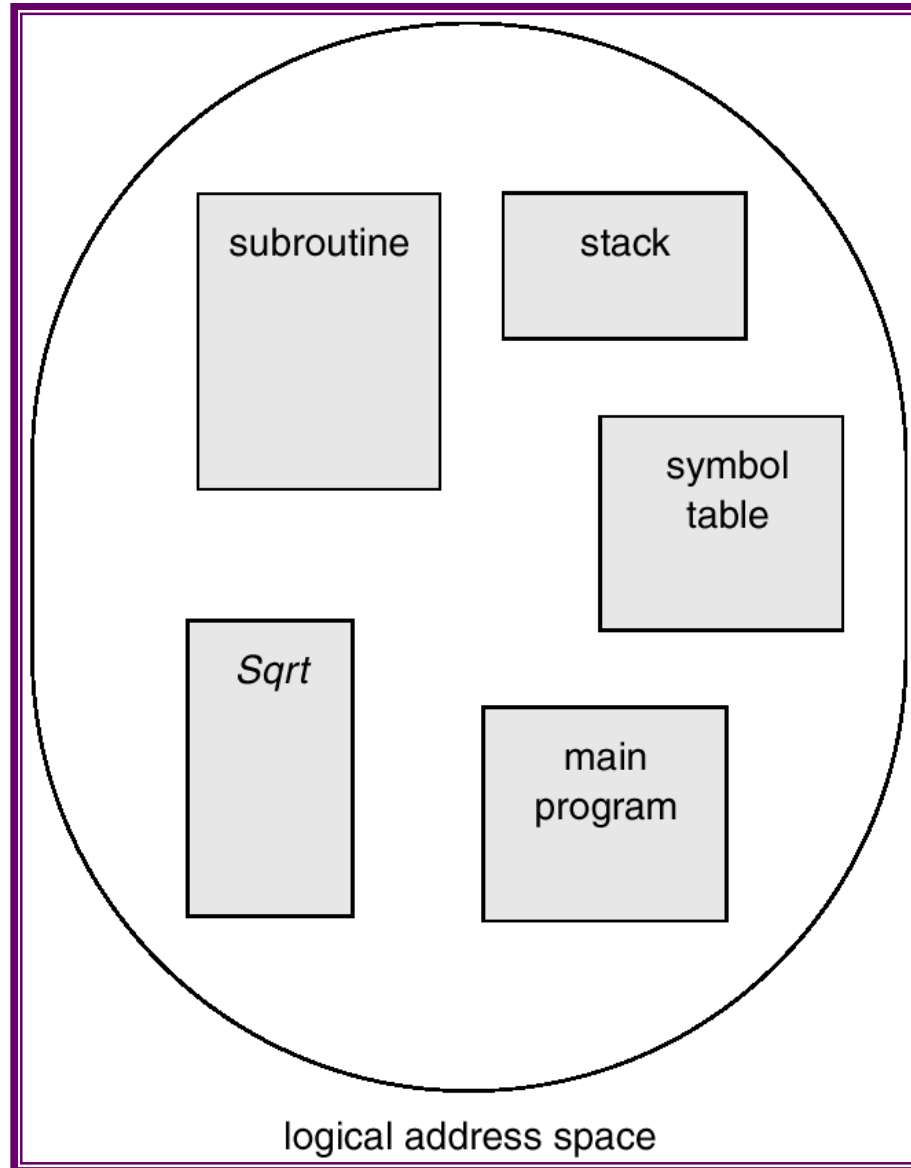
Shared Pages Example



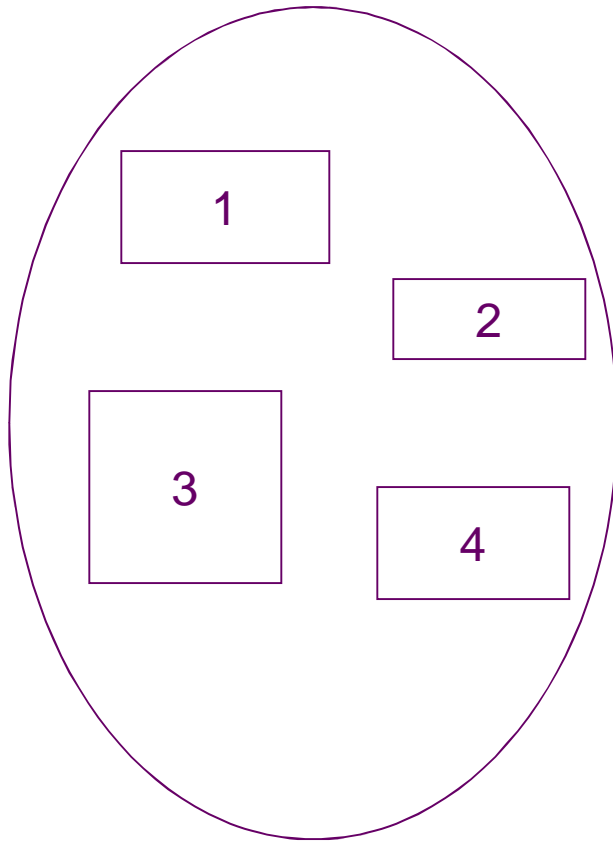
Segmentation

- Memory-management scheme that supports user view of memory.
- A program is a collection of segments. A segment is a logical unit such as:
 - main program,
 - procedure,
 - function,
 - method,
 - object,
 - local variables, global variables,
 - common block,
 - stack,
 - symbol table, arrays

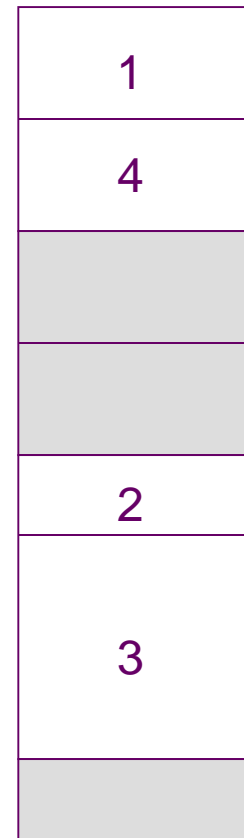
User's View of a Program



Logical View of Segmentation



user space



physical memory space

Segmentation Architecture

- Logical address consists of a two tuple:
 <segment-number, offset>,
- *Segment table* – maps two-dimensional physical addresses; each table entry has:
 - ☞ *base* – contains the starting physical address where the segments reside in memory.
 - ☞ *limit* – specifies the length of the segment.
- *Segment-table base register (STBR)* points to the segment table's location in memory.
- *Segment-table length register (STLR)* indicates number of segments used by a program;
 segment number s is legal if $s < \text{STLR}$.

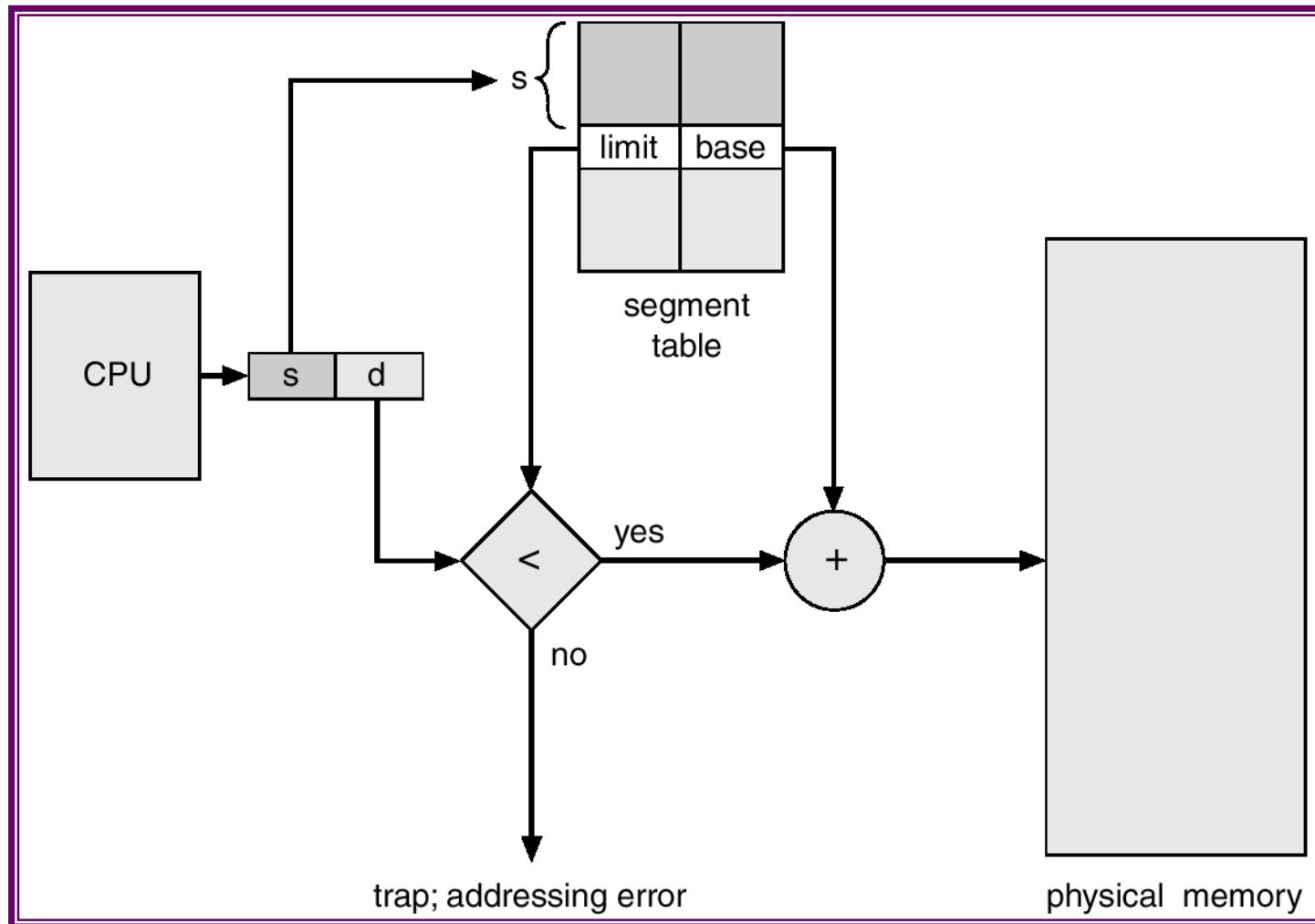
Segmentation Architecture (Cont.)

- Relocation.
 - ☞ dynamic
 - ☞ by segment table
- Sharing.
 - ☞ shared segments
 - ☞ same segment number
- Allocation.
 - ☞ first fit/best fit
 - ☞ external fragmentation

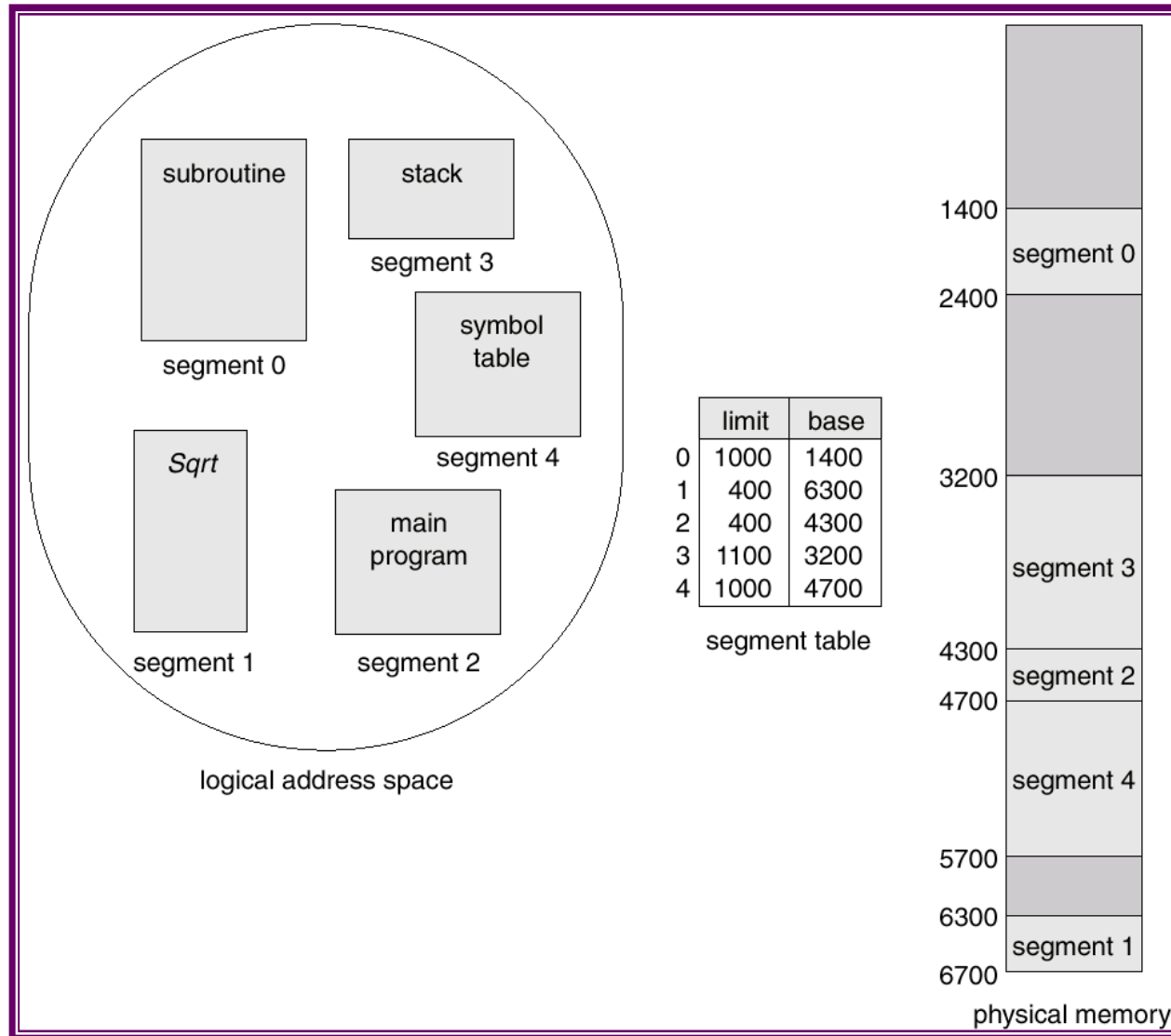
Segmentation Architecture (Cont.)

- Protection. With each entry in segment table associate:
 - ☞ validation bit = 0 \Rightarrow illegal segment
 - ☞ read/write/execute privileges
- Protection bits associated with segments; code sharing occurs at segment level.
- Since segments vary in length, memory allocation is a dynamic storage-allocation problem.
- A segmentation example is shown in the following diagram

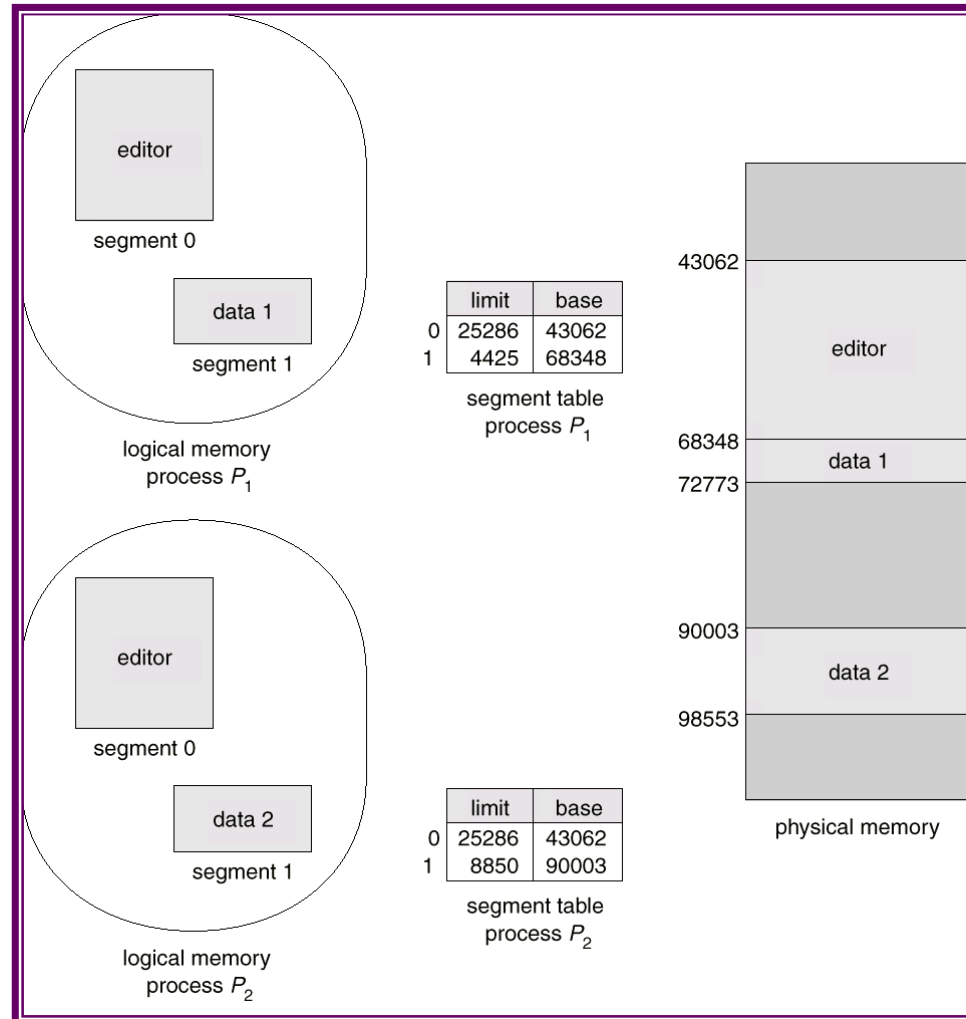
Segmentation Hardware



Example of Segmentation



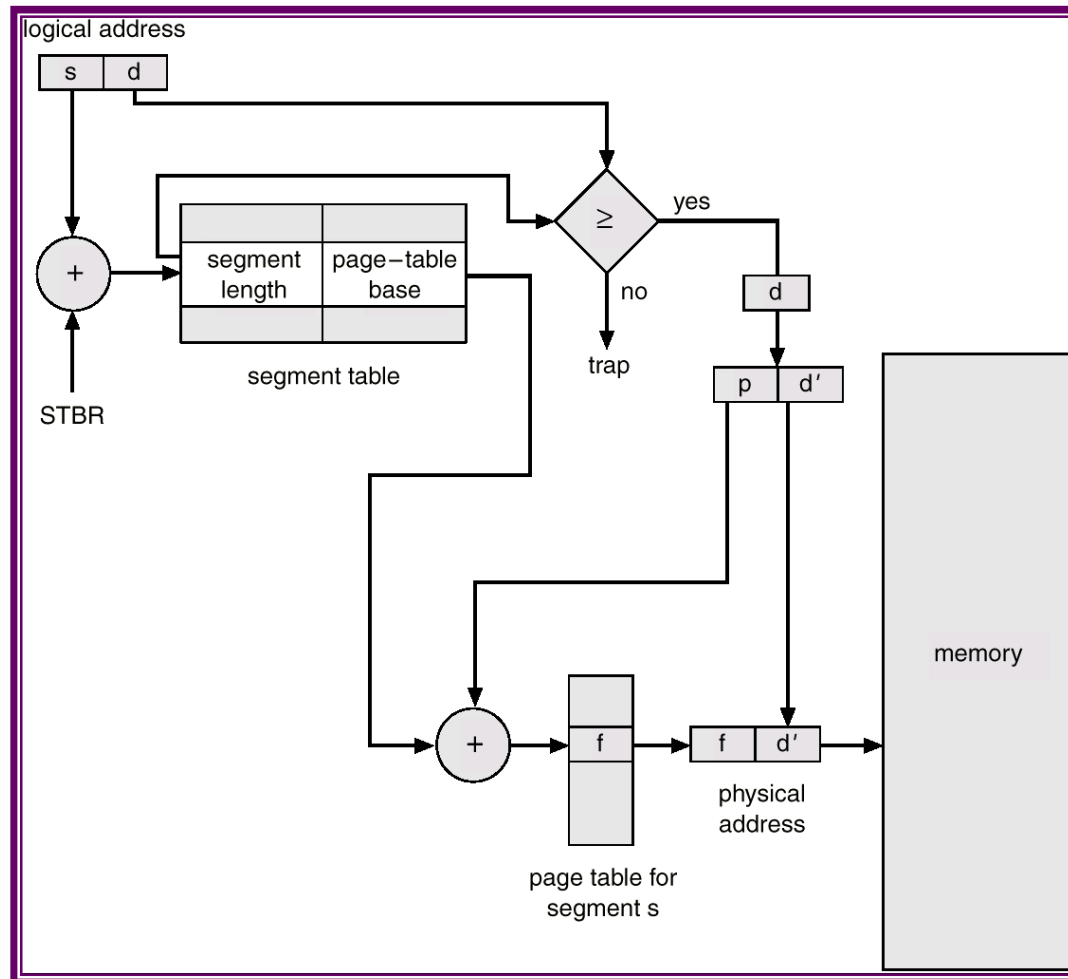
Sharing of Segments



Segmentation with Paging – MULTICS

- The MULTICS system solved problems of external fragmentation and lengthy search times by paging the segments.
- Solution differs from pure segmentation in that the segment-table entry contains not the base address of the segment, but rather the base address of a *page table* for this segment.

MULTICS Address Translation Scheme



Segmentation with Paging – Intel 386

- As shown in the following diagram, the Intel 386 uses segmentation with paging for memory management with a two-level paging scheme.

Intel 30386 Address Translation

